A Conditional Gradient Framework for Composite Convex Minimization with Applications to Semidefinite Programming

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Abstract

We propose a conditional gradient framework for a composite convex minimization template with broad applications. Our approach combines the notions of smoothing and homotopy under the CGM framework, and provably achieves the optimal $\mathcal{O}(1/\sqrt{k})$ convergence rate. We demonstrate that the same rate holds if the linear subproblems are solved approximately with additive or multiplicative error. Specific applications of the framework include the non-smooth minimization, semidefinite programming, and minimization with linear inclusion constraints over a compact domain. We provide numerical evidence to demonstrate the benefits of the new framework.

1 Introduction

The importance of convex optimization in machine learning has increased dramatically in the last decade due to the new theory in structured sparsity and rank minimization and statistical learning models like support vector machines. Indeed, a large class of learning formulations can be addressed by the following composite convex minimization template:

$$\min_{x \in \mathcal{X}} F(x) := f(x) + g(Ax),$$

(1.1)

where $\mathcal{X} \subset \mathbb{R}^n$ is compact (nonempty, bounded, closed) and its 0-dimensional faces (i.e., its vertices) are often called atoms. $f : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$ is a smooth proper closed convex function, $A \in \mathbb{R}^{d \times n}$, and $g : \mathbb{R}^d \to \mathbb{R} \cup \{+\infty\}$ is a proper closed convex function which is possibly non-smooth.

By using the powerful proximal gradient framework, the problems belonging to the template (1.1) can be solved nearly as efficiently as if they were fully smooth with fast convergence rates. By proximal (prox) operator, we mean the resolvent of the following optimization problem:

$$\text{prox}_g(v) = \arg \min_{x \in \mathbb{R}^d} g(x) + \frac{1}{2} \|x - v\|^2.$$

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These methods make use of the gradient of the smooth function $f$ along with the prox of the non-smooth function $g$, and are optimal in the sense that they match the iteration complexity lower-bounds.

Surprisingly, the proximal operator can impose an undesirable computational burden and even intractability on these gradient-based methods, such as the computation of a full singular value decomposition in the ambient dimension or the computation of proximal mapping for the latent group lasso [Jaggi, 2013]. Moreover, the linear mapping $A$ often complicates the computation of the prox itself, and require more sophisticated splitting or primal-dual methods.

As a result, the conditional gradient method (GCM, aka Frank-Wolfe method) has recently increased in popularity since it requires only a linear minimization oracle. By linear minimization oracle (lmo), we mean a resolvent of the following problem

$$\text{lmo}_\mathcal{X}(v) = \arg\min_{x \in \mathcal{X}} \langle x, v \rangle.$$  

The CGM features significantly reduced computational costs (e.g., when $\mathcal{X}$ is the spectrahedron), tractability (e.g., when $\mathcal{X}$ is a latent group lasso norm), and interpretability (e.g., they generate solutions as a combination of small number of extreme points of $\mathcal{X}$). The method, as shown in Algorithm 1 when $g(Ax) = 0$, is also simple to implement:

**Algorithm 1** CGM for smooth minimization

**Input:** $x_1 \in \mathcal{X}$

**for** $k = 1, 2, \ldots,$ **do**

$$\eta_k = \frac{2}{k+1},$$

$$s_k = \arg\min_{x \in \mathcal{X}} \langle \nabla f(x_k), x \rangle,$$

$$x_{k+1} = x_k + \eta_k (s_k - x_k).$$

**end for**

The method itself is optimal for this particular template since it achieves the iteration complexity lower-bound. Unfortunately, the CGM provably cannot handle the non-smooth $g(Ax)$ term in (1.1) (cf., Section 5.3 for a counter example by Nesterov [2017]).

When the non smooth part is an indicator function, one could take the intersection between $\mathcal{X}$ and the set represented by $g$. Unfortunately, even the lmo itself can be a difficult optimization problem depending on the structure of the domain. On many domains of interest that can be parametrized as a composition of simple sets, linear problems are infeasible [Richard et al., 2012, Yen et al., 2016].

In this paper, we propose a CGM framework for solving the composite problem (1.1) with rigorous convergence guarantees. Our approach retains the simplicity of projection free methods but allows to disentangle the complexity of the feasibility set in order to preserve the simplicity of the lmo.

Our method combines the ideas of smoothing [Nesterov, 2005] and homotopy under the CGM framework. Our study covers in particular the case where non-smooth part is the indicator function of a convex set. Similar ideas were proposed for the primal-dual subgradient method and the coordinate descent in [Tran-Dinh et al., 2017, Alacaoglu et al., 2017] via the projection onto $\mathcal{X}$.

Lan [2014] proposes a similar approach with the CGM for non-smooth problems, which is extended for the conditional gradient sliding framework in [Lan and Zhou, 2016, Lan et al., 2017]. Their analysis, however, is restricted by the assumption that the smoothed function is Lipschitz continuous. Consequently, it does not apply to the problems with affine inclusion constraints, limiting its applicability in machine learning (cf., Sections 5.5 and 5.6).

Our contributions can be summarized as follows:
We introduce a simple, easy to implement CGM framework for solving composite problem (1.1), and prove that it achieves the optimal $O(1/\sqrt{k})$ rate. To the best of our knowledge, our framework is the first CGM extension that can solve standard semidefinite programming formulation.

We analyze the convergence of our algorithm under inexact oracles with additive and multiplicative errors.

We present important special cases of our framework, including the non-smooth minimization, minimization with linear inclusion constraints, and minimization via splitting, along with the related work at each camp.

We present empirical evidence supporting our findings.

Roadmap. Section 2 recalls some basic notions and presents the preliminaries about the smoothing technique. In Section 3, we present CGM for composite convex minimization along with the convergence guarantees, and we extend these results for inexact oracle calls in Section 4. We describe some important special applications of our framework in Section 5. We provide empirical evidence supporting our theoretical findings in Section 6. Finally, Section 7 draws the conclusions with a discussion on the future work. Proofs and technical details are left to the appendix.

2 Notation & Preliminaries

We use $\| \cdot \|$ to denote the Euclidean norm for vectors and the spectral norm (a.k.a. Schatten $\infty$-norm) for linear mappings. We denote the Frobenius norm by $\| \cdot \|_F$, and the nuclear norm (a.k.a. Schatten 1-norm or trace norm) by $\| \cdot \|_{S_1}$. The notation $\langle \cdot, \cdot \rangle$ refers the Euclidean or Frobenius inner product. The symbol $\top$ denotes the adjoint of a linear map, and the symbol $\succeq$ denotes the semidefinite order. We denote the diameter of $\mathcal{X}$ by $D_{\mathcal{X}} = \max_{x_1, x_2 \in \mathcal{X}} \| x_1 - x_2 \|$.

Lipschitz continuity. We say that a function $g : \mathbb{R}^d \rightarrow \mathbb{R}$ is $L$-Lipschitz continuous if it satisfies

$$|g(x_1) - g(x_2)| \leq L \| x_1 - x_2 \|, \quad \forall x_1, x_2 \in \mathbb{R}^d.$$  

Smoothness. A differentiable function $f : \mathcal{X} \rightarrow \mathbb{R}$ is $L_f$-smooth if the gradient $\nabla f$ is $L_f$-Lipschitz continuous:

$$\| \nabla f(x_1) - \nabla f(x_2) \| \leq L_f \| x_1 - x_2 \|, \quad \forall x_1, x_2 \in \mathcal{X}.$$  

Fenchel conjugate & Smoothing. We consider the smooth approximation of a non-smooth term $g$ obtained using the technique described by Nesterov [2005] with the standard Euclidean proximity function $\frac{1}{2} \| \cdot \|^2$ and a smoothness parameter $\beta > 0$

$$g_\beta(z) = \max_{y \in \mathbb{R}^d} \langle z, y \rangle - g^*(y) - \frac{\beta}{2} \| y \|^2,$$

where $g^*$ denotes the Fenchel conjugate of $g$

$$g^*(x) = \sup_v \langle x, v \rangle - g(v).$$

Note that $g_\beta$ is convex and $\frac{1}{\beta}$-smooth. Throughout, we assume that $g$ is smoothing-friendly (cf., [Nesterov, 2005]), or a constraint indicator function.
**Solution set.** We denote an exact solution of (1.1) by $x^*$, and the set of all solutions by $X^*$. Throughout the paper, we assume that the solution set $X^*$ is nonempty.

Given an accuracy level $\epsilon > 0$, we call a point $x \in X$ as an $\epsilon$-solution of (1.1) if
\[
f(x) + g(Ax) - f^* - g^* \leq \epsilon, \tag{2.1}
\]
where we use the notation $f^* = f(x^*)$ and $g^* = g(Ax^*)$.

When $g$ is the indicator function of a set $K$, condition (2.1) is not well-defined for infeasible points. Hence, we refine our definition, and call a point $x \in X$ as an $\epsilon$-solution if
\[
f(x) - f^* \leq \epsilon, \quad \text{and} \quad \text{dist}(Ax, K) \leq \epsilon.
\]
Here, we call $f(x) - f^*$ as the objective residual and $\text{dist}(Ax, K)$ as the feasibility gap. We use the same $\epsilon$ for the objective residual and the feasibility gap, since the distinct choices can be handled by scaling $f$.

**Lagrange saddle point.** Suppose that $g$ is the indicator function of a convex set $K$. We assume that the Slater’s condition holds. By Slater’s condition, we mean
\[
\text{relint}(X \times K) \cap \{(x, r) \in \mathbb{R}^n \times \mathbb{R}^d : Ax = r\} \neq \emptyset,
\]
where relint stands for the relative interior. Denote the Lagrangian of problem (1.1) by
\[
\mathcal{L}(x, y) := f(x) + \langle y, Ax \rangle - g^*(y).
\]
We can formulate the primal and dual problems as follows:
\[
\sup_{y \in \mathbb{R}^d} \min_{x \in X} \mathcal{L}(x, y) \leq \min_{x \in X} \sup_{y \in \mathbb{R}^d} \mathcal{L}(x, y).
\]
We denote a solution of the dual problem by $y^*$.

### 3 Algorithm & Convergence

We design our main algorithm and present its convergence guarantees in this section.

Our method is based on the simple idea of combining smoothing and homotopy. Objective function $F$ in our problem template is non-smooth. We define the smooth approximation of $F$ with smoothness parameter $\beta > 0$ as
\[
F_\beta(x) = f(x) + g_\beta(Ax).
\]
Note that $F_\beta$ is $(L_f + \|A\|^2/\beta)$-smooth.

The algorithm takes a conditional gradient step with respect to the smooth approximation $F_\beta_k$ at iteration $k$, where $\beta_k$ is gradually decreased towards 0.

Let us denote by $y^*_{\beta_k}$
\[
y^*_{\beta_k}(Ax) = \arg \max_{y \in \mathbb{R}^d} \langle Ax, y \rangle - g^*(y) - \frac{\beta_k}{2} \|y\|^2
\]
\[
= \text{prox}_{\beta_k^{-1}g}(\beta_k^{-1}Ax)
\]
\[
= \frac{1}{\beta_k}(Ax - \text{prox}_{\beta_k g}(Ax)),
\]

where the last equality is due to the Moreau decomposition. Then, we can compute the gradient of $F_{\beta_k}$ as

$$
\nabla F_{\beta_k}(x) = \nabla f(x) + A^\top y_{\beta_k}(Ax)
$$

$$
= \nabla f(x) + \frac{1}{\beta_k} A^\top (Ax - \text{prox}_{\beta_k g}(Ax)).
$$

Based on this formulation, we present our CGM framework for composite convex minimization template (1.1) in Algorithm 2. The choice of $\beta_k$ comes from the convergence analysis, which can be found in the supplements.

**Algorithm 2** CGM for composite problems

**Input:** $x_1 \in \mathcal{X}$, $\beta_0 > 0$

for $k = 1, 2, \ldots$, do

$\eta_k = \frac{2}{k+1}$, and $\beta_k = \frac{\beta_0}{\sqrt{k+1}}$

$\nu_k = \beta_k \nabla f(x_k) + A^\top (Ax_k - \text{prox}_{\beta_k g}(Ax_k))$

$s_k = \arg \min_{x \in \mathcal{X}} \langle \nu_k, x \rangle$

$x_{k+1} = x_k + \eta_k (s_k - x_k)$

end for

**Theorem 3.1.** The sequence $x_k$ generated by Algorithm 2 satisfies the following bound:

$$
F_{\beta_k}(x_{k+1}) - F^* \leq 2D_X^2 \left( \frac{L_f}{k+1} + \frac{||A||^2}{\beta_0 \sqrt{k+1}} \right).
$$

Theorem 3.1 does not directly certify the convergence of $x_k$ to the solution, since the bound is on the smoothed gap $F_{\beta_k}(x_k) - F^*$. However, it is a milestone to prove the convergence guarantees in Theorems 3.2 and 3.3.

**Theorem 3.2.** Assume that $g : \mathbb{R}^d \to \mathbb{R}$ is $L_g$-Lipschitz continuous. Then, the sequence $x_k$ generated by Algorithm 2 satisfies the following convergence bound:

$$
F(x_k) - F^* \leq 2D_X^2 \left( \frac{L_f}{k} + \frac{||A||^2}{\beta_0 \sqrt{k}} \right) + \frac{\beta_0 L_g^2}{2\sqrt{k}}.
$$

Furthermore, if the constants $D_X$, $||A||$ and $L_g$ are known or easy to approximate, we can choose $\beta_0 = 2D_X ||A||/L_g$ to get the following convergence rate:

$$
F(x_k) - F^* \leq \frac{2D_X^2 L_f}{k} + \frac{2D_X ||A|| L_g}{\sqrt{k}}.
$$

**Theorem 3.3.** Assume that $g : \mathbb{R}^d \to \mathbb{R}$ is the indicator function of a simple convex set $\mathcal{K}$. Then, the sequence $x_k$ generated by Algorithm 2 satisfies:

$$
f(x_k) - f^* \geq -||y^*|| \text{ dist}(Ax_k, \mathcal{K})
$$

$$
f(x_k) - f^* \leq 2D_X^2 \left( \frac{L_f}{k} + \frac{||A||^2}{\beta_0 \sqrt{k}} \right)
$$

$$
\text{dist}(Ax_k, \mathcal{K}) \leq \frac{2\beta_0}{\sqrt{k}} \left( ||y^*|| + D_X \sqrt{\frac{C_0}{\beta_0}} \right)
$$

where $C_0 = L_f + ||A||^2 / \beta_0$. 


Remark 3.4. Similar to the CGM for smooth minimization, we can consider variants of Algorithm 2 with line-search and fully corrective updates (cf., [Jaggi, 2013]). Theorems 3.1 to 3.3, as well as their extensions for inexact oracle calls in Section 4, still hold for the variants of Algorithm 2 with line-search (which replaces the step size by \( \eta_k = \min_{\eta \in [0,1]} F_\beta_k (x_k + \eta (s_k - x_k)) \)) and fully corrective updates (which replaces the last step by \( x_{k+1} = \arg \min_{x \in \text{conv}(s_1, \ldots, s_k)} F_\beta_k (x) \)).

4 Convergence with Inexact Oracles

Finding an exact solution of the lmo can be expensive in practice, especially when it involves a matrix factorization as in the SDP examples. On the other hand, approximate solutions can be much more efficient.

Different notions of inexact lmo are already explored in the Frank-Wolfe and greedy optimization frameworks, cf., [Lacoste-Julien et al., 2013, Locatello et al., 2017a,b]. We revisit the notions of additive and multiplicative errors which we adapt here for our setting.

4.1 Inexact Oracle with Additive Error

At iteration \( k \), for the given direction \( v_k \), we assume that the approximate lmo returns an element \( \tilde{s}_k \in X \) such that:

\[
\langle v_k, \tilde{s}_k \rangle \leq \langle v_k, s_k \rangle + \delta \eta_k 2D_X^2 \left( L_f + \frac{\|A\|^2}{\beta_k} \right)
\]  (4.1)

for some \( \delta > 0 \). Note that as in [Jaggi, 2013], we require the accuracy of lmo to increase as the algorithm progresses.

Replacing the exact lmo with the approximate oracles of the form (4.1) in Algorithm 2, we get the convergence guarantees in Theorems 4.1 to 4.3.

Theorem 4.1. The sequence \( x_k \) generated by Algorithm 2 with approximate lmo (4.1) satisfies:

\[
F_\beta_k (x_{k+1}) - F^* \leq 2D_X^2 \left( \frac{L_f}{k+1} + \frac{\|A\|^2}{\beta_k \sqrt{k+1}} \right) (1 + \delta).
\]

Theorem 4.2. Assume that \( g \) is \( L_g \)-Lipschitz continuous. Then, the sequence \( x_k \) generated by Algorithm 2 with approximate lmo (4.1) satisfies:

\[
F(x_k) - F^* \leq 2D_X^2 \left( \frac{L_f}{k} + \frac{\|A\|^2}{\beta_k \sqrt{k}} \right) (1 + \delta) + \frac{\beta_0 L_g^2}{2\sqrt{k}}.
\]

We can optimize \( \beta_0 \) from this bound if \( \delta \) is known.

Theorem 4.3. Assume that \( g \) is the indicator function of a simple convex set \( K \). Then, the sequence \( x_k \) generated by Algorithm 2 with approximate lmo (4.1) satisfies:

\[
f(x_k) - f^* \geq -\|y^*\| \text{dist}(Ax_k, K)
\]

\[
f(x_k) - f^* \leq 2D_X^2 \left( \frac{L_f}{k} + \frac{\|A\|^2}{\beta_k \sqrt{k}} \right) (1 + \delta)
\]

\[
\text{dist}(Ax_k, K) \leq \frac{2\beta_0}{\sqrt{k}} \left( \|y^*\| + D_X \sqrt{\frac{C_0}{\beta_0} (1 + \delta)} \right)
\]

where \( C_0 = L_f + \|A\|^2/\beta_0 \).
4.2 Inexact Oracle with Multiplicative Error

We consider the multiplicative inexact oracle:

\[ \langle v_k, \tilde{s}_k - x_k \rangle \leq \delta \langle v_k, s_k - x_k \rangle \]  

(4.2)

where \( \delta \in (0, 1] \). Replacing the exact lmo with the approximate oracles of the form (4.2) in Algorithm 2, we get the convergence guarantees in Theorems 4.4 to 4.6.

**Theorem 4.4.** The sequence \( x_k \) generated by Algorithm 2 with approximate lmo of the form (4.2), and modifying \( \eta_k = \frac{2}{\delta(k-1)+2} \) and \( \beta_k = \frac{\beta_0}{\sqrt{\delta k+1}} \) satisfies:

\[ F_{\beta_k}(x_{k+1}) - F^* \leq \frac{2}{\delta} \left( \frac{D_A^2 L_f + \delta E}{\delta k + 2} + \frac{D_A^2 \|A\|^2}{\beta_0 \sqrt{\delta k + 2}} \right) \]

where \( E = F(x_1) - F^* \).

**Theorem 4.5.** Assume that \( g \) is \( L_g \)-Lipschitz continuous. Then, the sequence \( x_k \) generated by Algorithm 2 with approximate lmo (4.2), and modifying \( \eta_k = \frac{2}{\delta(k-1)+2} \) and \( \beta_k = \frac{\beta_0}{\sqrt{\delta k+1}} \) satisfies:

\[ F(x_k) - F^* \leq \frac{2}{\delta} \left( \frac{D_A^2 L_f + \delta E}{\delta k + 1} + \frac{D_A^2 \|A\|^2}{\beta_0 \sqrt{\delta k + 1}} \right) + \frac{\beta_0 L_g^2}{2 \sqrt{\delta k + 1}} \]

where \( E = F(x_1) - F^* \). We can optimize \( \beta_0 \) from this bound if \( \delta \) is known.

**Theorem 4.6.** Assume that \( g \) is the indicator function of a simple convex set \( K \). Then, the sequence \( x_k \) generated by Algorithm 2 with approximate lmo (4.2), and modifying \( \eta_k = \frac{2}{\delta(k-1)+2} \) and \( \beta_k = \frac{\beta_0}{\sqrt{\delta k+1}} \) satisfies:

\[ f(x_k) - f^* \geq - \|y^*\| \text{dist}(A x_k, K) \]

\[ f(x_k) - f^* \leq \frac{2}{\delta} \left( \frac{D_A^2 L_f + \delta E}{\delta k + 1} + \frac{D_A^2 \|A\|^2}{\beta_0 \sqrt{\delta k + 1}} \right) \]

\[ \text{dist}(A x_k, K) \leq \frac{2 \beta_0}{\sqrt{\delta k + 1}} \left( \|y^*\| + \sqrt{\frac{D_A^2 E_0 + \delta E}{\beta_0 \delta}} \right) \]

where \( E = F(x_1) - F^* \) and \( E_0 = L_f + \|A\|^2/\beta_0 \).

5 Applications & Related Work

The CGM is proposed for the first time in the seminal work of Frank and Wolfe [1956] for solving smooth convex optimization on a polytope. It is then progressively generalized for more general settings in [Levitin and Polyak, 1966, Dunn and Harshbarger, 1978, Dunn, 1979, 1980]. Nevertheless, with the introduction of the fast gradient methods with \( O(1/k^2) \) rate by Nesterov [1987], the development of CGM-type methods entered into a stagnation period.

The recent developments in machine learning applications with vast data brought the scalability of the first order methods under scrutiny. As a result, there has been a renewed interest in the CGM in the last decade. Hence, we compare our framework with the recent developments of CGM literature in different camps of problem templates below.
5.1 Smooth Problems

The CGM is extended for the smooth convex minimization over the simplex by Clarkson [2010], for the spectrahedron by Hazan [2008], and for an arbitrary compact convex set by Jaggi [2013]. Online, stochastic and block coordinate variants of CGM are introduced by Hazan and Kale [2012], Hazan and Luo [2016] and Lacoste-Julien et al. [2013] respectively.

When applied to smooth problems, Algorithm 2 is equivalent to the classical CGM, and Theorem 3.2 recovers the known optimal $O(1/k)$ convergence rate. We refer to [Jaggi, 2013] for a review of applications of this template.

It needs to be mentioned that Nesterov [2017] relaxes the smoothness assumption showing that the CGM converges for weakly-smooth objectives (i.e., with Hölder continuous gradients of order $\nu \in (0, 1)$).

5.2 Regularized Problems

The CGM for composite problems is considered recently by Nesterov [2017] and Xu [2017]. A similar but slightly different template, where $\mathcal{X}$ and $g$ are assumed to be a closed convex cone and a norm respectively, is also studied by [Harchaoui et al., 2015]. However, these works are based on the resolvents of a modified oracle,

$$\arg\min_{x \in \mathcal{X}} \langle x, v \rangle + g(Ax),$$

which can be expensive, unless $\mathcal{X} \equiv \mathbb{R}^n$, or $g = 0$.

Algorithm 2 applies to the problem template (1.1) by leveraging prox of the regularizer and lmo of the domain independently. This allows us to consider additional sparsity, group sparsity and structured sparsity promoting regularizations, elastic-net regularization, total variation regularization and many others under the CGM framework.

Semi-proximal mirror-prox proposed by [He and Harchaoui, 2015] also based on the smoothing technique, yet the motivation is fundamentally different. This method considers the regularizers for which the prox is difficult to compute, but can be approximated via CGM.

5.3 Non-Smooth Problems

Template (1.1) covers the non-smooth convex minimization template as a special case:

$$\min_{x \in \mathcal{X}} g(Ax). \tag{5.1}$$

Unfortunately, the classical CGM (Algorithm 1) cannot handle the non-smooth minimization, as shown by Nesterov [2017] with the following counter-example.

**Example.** Let $\mathcal{X}$ be the unit Euclidean norm ball in $\mathbb{R}^2$, and $g(x) = \max\{x_1, x_2\}$. Clearly, $x^* = [\sqrt{2}, -\sqrt{2}]^T$. Choose an initial point $x_0 \neq x^*$. We can use an oracle that returns a subgradient $\nabla f(x) \in \begin{bmatrix} 1 & 0 \end{bmatrix}^T$ at any point $x \in \mathcal{X}$. Therefore, lmo returns $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ or $\begin{bmatrix} 0 \\ 1 \end{bmatrix}$ at each iteration, and $x_k$ belongs to the convex hull of $\{x_0, \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \end{bmatrix}\}$ which does not contain the solution.

Our framework escapes such issues by leveraging prox of the objective function $g$. In this pathological example, prox$_g$ corresponds to the projection onto the simplex. Often times the cost of prox$_g$ is negligible in comparison to the cost of lmo$_\mathcal{X}$ (cf., Section 6.2 for a robust PCA example).

Assume that $g : \mathbb{R}^d \to \mathbb{R}$ is $L_g$-Lipschitz continuous. As a consequence of Theorem 3.2, Algorithm 2 for solving (5.1) by choosing $\beta_0 = 2D_\mathcal{X}\|A\|/L_g$ satisfies

$$g(Ax_k) - g^* \leq \frac{2D_\mathcal{X}\|A\|L_g}{\sqrt{k}}.$$
Figure 1: The classical CGM (left) and our framework (right) for the pathological example, starting from $\frac{1}{10}$.

We recover the method proposed by Lan [2014] in this specific setting. Lan [2014] shows that this rate is optimal for algorithms approximating the solution of (5.1) as a convex combination of resolvents of lmo.

Our analysis with inexact oracles in this setting is new. In stark contrast to the smooth case, where the additive error should decrease by $O(1/k)$ rate, definition (4.1) implies that we can preserve the convergence rate in the non-smooth case if the additive error is $O(1/\sqrt{k})$.

5.4 Minimax Problems

We consider the minimax problems of the following form:

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \mathcal{L}(Ax, y)$$

where $\mathcal{L}$ is a smooth convex-concave function, i.e., $\mathcal{L}(\cdot, y)$ is convex $\forall y \in \mathcal{Y}$ and $\mathcal{L}(Ax, \cdot)$ is concave $\forall x \in \mathcal{X}$. Note that this formulation is a special instance of (5.1) with $g(Ax) = \max_{y \in \mathcal{Y}} \mathcal{L}(Ax, y)$. Consequently, we can apply Algorithm 2 if $\text{prox}_g$ is tractable.

When $\mathcal{Y}$ admits an efficient projection oracle, $\text{prox}_g$ is also efficient for bilinear saddle point problems $\mathcal{L}(Ax, y) = \langle Ax, y \rangle$. By Moreau decomposition, we have

$$\text{prox}_y(Ax_k) = Ax_k - \text{proj}_y(Ax_k),$$

hence $v_k$ takes the form

$$v_k = \beta_k \nabla f(x_k) + A^\top \text{proj}_y(Ax_k).$$

Gidel et al. [2017] proposes a CGM variant for the smooth convex-concave saddle point problems. This method process both $x$ and $y$ via the lmo, and hence it also requires $\mathcal{Y}$ to be bounded. Our method, on the other hand, is more suitable when $\text{proj}_y$ is easy.

Bilinear saddle point problem covers the maximum margin estimation of structured output models [Taskar et al., 2006] and minimax games [Von Neumann and Morgenstern, 1944].

In particular, it also covers an important semidefinite programming formulation [Garber and Hazan, 2016], where $\mathcal{X}$ is a spectrahedron and $\mathcal{Y}$ is the simplex. Our framework fits perfectly here since the projection onto the simplex can be computed efficiently. Here, we defer the derivation of the extension of our framework with the entropy Bregman smoothing for future.
5.5 Problems with Affine Constraints

Algorithm 2 also applies to smooth convex minimization problems with affine constraints over a convex compact set:

$$\min_{x \in \mathcal{X}} f(x) \quad \text{subject to} \quad Ax = b,$$

(5.2)

by setting $g(Ax)$ in (1.1) as indicator function of set $\{b\}$, where $b \in \mathbb{R}^d$ is a known vector.

Since the prox operator of the indicator function of a convex set is the projection, $v_k$ in Algorithm 2 becomes

$$v_k = \beta_k \nabla f(x_k) + A^T (Ax_k - b).$$

A relevant approach to our framework in this setting is the universal primal-dual gradient method (UPD) proposed by Yurtsever et al. [2015]. The UPD method takes advantage of Fenchel-type oracles, which can be thought as a generalization of the lmo. The method is based on an inexact line-search technique in the dual, and recovers the primal variable via averaging.

Unfortunately, UPD iterations explicitly depend on the target accuracy level $\epsilon$, which is difficult to tune since it requires rough knowledge of the optimal value. Moreover, the method converges only up to $\epsilon$-suboptimality. There is no known analysis with inexact oracle calls for UPD, and errors in function evaluation can cause the algorithm to get stuck in the line-search procedure.

We can generalize (5.2) for the problems with affine inclusion constraints:

$$\min_{x \in \mathcal{X}} f(x) \quad \text{subject to} \quad Ax - b \in \mathcal{K},$$

(5.3)

where $\mathcal{K}$ is a simple closed convex set. In this case, $v_k$ takes the following form:

$$v_k = \beta_k \nabla f(x_k) + A^T (Ax_k - b - \text{proj}_{\mathcal{K}}(Ax_k - b)).$$

We implicitly assume that $\text{proj}_{\mathcal{K}}$ is tractable. We can use a splitting framework when it is computationally more advantageous to use lmo$\mathcal{K}$ instead (cf. Section 5.6).

This template covers the standard semidefinite programming in particular. Applications include clustering [Peng and Wei, 2007], optimal power-flow [Lavaei and Low, 2012], sparse PCA [d’Aspremont et al., 2007], kernel learning [Lanckriet et al., 2004], blind deconvolution [Ahmed et al., 2014], community detection [Bandeira et al., 2016], etc. Besides machine learning applications, this formulation has a crucial role in the convex relaxation of combinatorial problems.

A significant example is the problems over the doubly nonnegative cone (i.e., the intersection of the positive semidefinite cone and the positive orthant) with a bounded trace norm [Yoshise and Matsukawa, 2010]. Note that the lmo over this domain can be costly since the lmo can require full dimensional updates [Hamilton-Jester and Li, 1996, Locatello et al., 2017b].

Our framework can handle these problems ensuring the positive semidefiniteness by lmo$\mathcal{X}$, and can still ensure the convergence to the first orthant via $\text{proj}_{\mathcal{K}}$.

To the best of our knowledge, our framework is the first CGM extension that can handle affine constraints.

5.6 Minimization via Splitting

We can take advantage of a splitting framework since we can handle affine constraints. This lets us to disentangle the complexity of the feasibility set.
Consider the following optimization template:

$$\min_{x \in X_1 \cap X_2} f(x) + g(Ax)$$

subject to

$$Bx - b \in K$$

$$Cx - c \in S$$

where $X_1$ and $X_2 \subset \mathbb{R}^n$ are two convex compact sets, $A, B, C$ are known matrices and $b, c$ are given vectors.

Suppose that

- \(\text{lmo}_{X_1}\) and \(\text{lmo}_{X_2}\) are easy to compute, but not \(\text{lmo}_{X_1 \cap X_2}\)
- \(\text{prox}_g\) is easy to compute
- \(K\) is a simple convex set and \(\text{proj}_K\) is efficient
- \(S\) is a convex compact set with an efficient lmo.

We can reformulate this problem introducing slack variables $\xi \in X_2$ and $\psi \in S$ as follows:

$$\min_{x \in X_1, \xi \in X_2, \psi \in S} f(x) + g(Ax)$$

subject to

$$Bx - b \in K$$

$$Cx - c = \psi, \ x = \xi.$$ 

This formulation is in the form of (5.2) with respect to the variable $(x, \xi, \psi) \in X_1 \times X_2 \times S$. Therefore, we can apply Algorithm 2. It is easy to see that Algorithm 2 leverages \(\text{lmo}_{X_1}\), \(\text{lmo}_{X_2}\), \(\text{lmo}_S\), \(\text{prox}_g\) and \(\text{proj}_K\) separately.

We can generalize this approach in a straightforward way for problems with an arbitrary finite number of non-smooth terms:

$$\min_{x \in \bigcap_i X_i} f(x) + \sum_j g_j(A_jx)$$

subject to

$$B_\ell x - b_\ell \in K_\ell$$

$$C_mx - c_m \in S_m.$$ 

6 Numerical Experiments

This section presents numerical experiments supporting our theoretical findings in clustering and robust PCA examples. The non-smooth parts in the chosen examples consist of indicator functions, for which the dual domain is unbounded. Hence, to the best of our knowledge, other CGM variants in the literature are not applicable.

6.1 Clustering the MNIST dataset

We consider the model-free $k$-means clustering based on the semidefinite relaxation of Peng and Wei [2007]:

$$\min_{X \in \mathcal{X}} \langle D, X \rangle \quad \text{subject to} \quad \underbrace{X1 = 1}_{g}, \quad X \succeq 0, \quad \text{tr}(X) \leq \rho.$$ 

where $\mathcal{X} = \{X \in \mathbb{R}^{n \times n} : \ X \succeq 0, \ \text{tr}(X) \leq \rho\}$ is the set of positive semidefinite matrices with a bounded trace norm, and $D \in \mathbb{R}^{n \times n}$ is the Euclidean distance matrix.

We use the test setup described and published online by Mixon et al. [2017], which can be briefly described as follows: First the meaningful features from MNIST dataset [LeCun and Cortes], which
Figure 2: Clustering MNIST dataset: Convergence of our framework in function value and the feasibility gap. Red dashed line on the misclassification plot represents the value reported by Mixon et al. [2017].

consists of 28 × 28 grayscale images that can be stacked as 784 × 1 vectors, are extracted using a one-layer neural network. This amounts to finding a weight matrix $W \in \mathbb{R}^{784 \times 10}$ and a bias vector $b \in \mathbb{R}^{10}$. Then, the trained neural network is applied to the first 1000 elements of the test set, which gives the probability vectors for these 1000 test points, where each entry represents the probability of being each digit.

Mixon et al. [2017] runs a relax-and-round algorithm which solves (6.1) by SDPNAL+ [Yang et al., 2015] followed by a rounding scheme (see Section 5 of [Mixon et al., 2017] for details), and compares the results against MATLAB’s built-in $k$-means++ implementation. Relax-and-round method is reported to achieve a misclassification rate of 0.0971. This rate matches with the all-time best rate for $k$-means++ after 100 different runs with random initializations.

For this experiment, we solve (6.1) by using Algorithm 2. Then, we cluster data using the same rounding scheme as [Mixon et al., 2017]. We initialize our method from the matrix of all zeros, and we choose $\beta_0 = 1$. We solve the lmo using the built-in MATLAB `eigs` function with the tolerance parameter $10^{-9}$.

We present the results of this experiment in Figure 2. We observe empirical $O(1/\sqrt{k})$ rate both in the objective residual and the feasibility gap. Surprisingly, the method achieves the best test error around 1000 iterations achieving the misclassification rate of 0.0914. This improves the value reported in [Mixon et al., 2017] by 5.8%.

This example demonstrates that the slow convergence rate is not a major problem in many machine learning problems, since a low accuracy solution can generalize as well as the optimal point in terms of the test error, if not better.

6.2 Robust PCA

Suppose that we are given a large matrix that can be decomposed as the summation of a low-rank and a sparse (in some representation) matrix. Robust PCA aims to recover these components accurately. Robust PCA has many applications in machine learning and data science, such as collaborative filtering, system identification, genotype imputation, etc. Here, we focus on an image decomposition problem so that we can visualize the decomposition error results.

Our setting is similar to the setup described in [Zeng and So, 2018]. We consider a scaled grayscale photograph with pattern from [Liu et al., 2013], and we assume that we only have access to an occluded image. Moreover, the image is contaminated by salt and pepper noise of density 1/10. We seek to approximate the original from this noisy image.

This is essentially a matrix completion problem, and most of the scalable techniques rely on the
Gaussian noise model. Note however the corresponding least-squares formulation is a good model against outliers:

$$\min_{X \in \mathcal{X}} \frac{1}{2} \| A(X) - b \|^2 \quad \text{subject to} \quad 0 \leq X \leq 1,$$

where $\mathcal{X} = \{ X \in \mathbb{R}^{n \times n} : \| X \|_{S_1} \leq \rho \}$ is a scaled nuclear norm ball, and $A : \mathbb{R}^{n \times n} \to \mathbb{R}^d$ is the sampling operator.

Our framework also covers the following least absolute deviations formulation which is known to be more robust:

$$\min_{X \in \mathcal{X}} \| A(X) - b \|_1 \quad \text{subject to} \quad 0 \leq X \leq 1.$$

We solve both formulations with our framework, starting from all zero matrix, running 1000 iterations, and assuming that we know the true nuclear norm of the original image. We choose $\beta_0 = 1$ in both cases.

This experiment demonstrates the implications of the flexibility of our framework in a simple machine learning setup. We compile the results in Figure 3, where the non-smooth formulation recovers a better approximation with 5dB higher peak signal to noise ratio (PSNR) and 0.27 higher structural similarity index (SSIM). Evaluation of PSNR and SSIM vs iteration counter are shown in Figure 4.

7 Conclusion

We presented a CGM framework for the composite convex minimization template, that provably achieves the optimal rate. This rate also holds under approximate oracle calls with additive or multiplicative errors.
Figure 4: Evaluation of PSNR and SSIM vs iteration counter.

Apart from its generalizations for various templates, there has been many attempts to improve the convergence rate, the arithmetic and the storage cost, or the proof techniques of the CGM under some specific settings, cf. [Dunn, 1979, Guélat and Marcotte, 1986, Beck, 2004, Garber and Hazan, 2015, Lacoste-Julien and Jaggi, 2015, Odor et al., 2016, Freund and Grigas, 2016, Yurtsever et al., 2017] and the references therein.

Many of these techniques can be adapted in our framework, since we preserve the key features of the CGM, such as the reduced costs and the atomic representations. The only seeming drawback is the loss of affine invariance in the analysis, left for future, which is fundamentally challenging due to smoothing technique.

References


Appendix

A1 Preliminaries

The following properties of smoothing are key to derive the convergence rate of our algorithm.

Let $g : \mathbb{R}^d \to \mathbb{R} \cap \{+\infty\}$ be a proper, closed and convex function, and denote its smooth approximation by

$$g_\beta(z) = \max_{y \in \mathbb{R}^d} \langle z, y \rangle - g^*(y) - \frac{\beta}{2} \|y\|^2$$

where $g^*$ represents the Fenchel conjugate of $g$ and $\beta > 0$ is the smoothing parameter. Then, $g_\beta$ is convex and $\frac{1}{\beta}$-smooth. Let us denote the unique maximizer of this concave problem by

$$y_\beta^*(z) = \arg \max_{y \in \mathbb{R}^d} \langle z, y \rangle - g^*(y) - \frac{\beta}{2} \|y\|^2$$

where $g^*$ represents the Fenchel conjugate of $g$ and $\beta > 0$ is the smoothing parameter. Then, $g_\beta$ is convex and $\frac{1}{\beta}$-smooth. Let us denote the unique maximizer of this concave problem by

$$y_\beta^*(z) = \arg \max_{y \in \mathbb{R}^d} \langle z, y \rangle - g^*(y) - \frac{\beta}{2} \|y\|^2$$

where the last equality is known as the Moreau decomposition. Then, the followings hold for any $z_1, z_2 \in \mathbb{R}^d$, and any $\beta, \gamma > 0$

$$g_\beta(z_1) \geq g_\beta(z_2) + \langle \nabla g_\beta(z_2), z_1 - z_2 \rangle + \frac{\beta}{2} \|y_\beta^*(z_2) - y_\beta^*(z_1)\|^2 \quad (7.1)$$

$$g(z_1) \geq g_\beta(z_2) + \langle \nabla g_\beta(z_2), z_1 - z_2 \rangle + \frac{\beta}{2} \|y_\beta^*(z_2)\|^2 \quad (7.2)$$

$$g_\beta(z_1) \leq g_\gamma(z_1) + \frac{\gamma - \beta}{2} \|y_\beta^*(z_1)\|^2 \quad (7.3)$$

Proofs can be found in Lemma 10 from [Tran-Dinh et al., 2017].

Suppose that $g$ is $L_g$-Lipschitz continuous. Then, for any $\beta > 0$ and any $z \in \mathbb{R}^d$, the following bound holds:

$$g_\beta(z) \leq g(z) \leq g_\beta(z) + \frac{\beta}{2} L_g^2 \quad (7.4)$$

Proof follows from equation (2.7) in [Nesterov, 2005] with a remark on the duality between Lipshitzness and bounded support (cf. Lemma 5 in [Dünner et al., 2016]).

A2 Convergence analysis

This section presents the proof of our convergence results. We skip proofs of Theorems 3.1 to 3.3 since we can get these results as a special case by setting $\delta = 0$ in Theorems 4.1 to 4.3.
Proof of Theorem 4.1

First, we use the smoothness of $F_{\beta_k}$ to upper bound the progress. Note that $F_{\beta_k}$ is $(L_f + \|A\|^2/\beta_k)$-smooth.

$$F_{\beta_k}(x_{k+1}) \leq F_{\beta_k}(x_k) + \eta_k \langle \nabla F_{\beta_k}(x_k), \tilde{s}_k - x_k \rangle + \frac{\eta_k^2}{2} \|\tilde{s}_k - x_k\|^2 (L_f + \frac{\|A\|^2}{\beta_k})$$

$$\leq F_{\beta_k}(x_k) + \eta_k \langle \nabla F_{\beta_k}(x_k), \tilde{s}_k - x_k \rangle + \frac{\eta_k^2}{2} D^2_{\mathcal{X}}(L_f + \frac{\|A\|^2}{\beta_k}), \tag{7.5}$$

where $\tilde{s}_k$ denotes the atom selected by the inexact linear minimization oracle, and the second inequality follows since $\tilde{s}_k \in \mathcal{X}$.

By definition of inexact oracle (4.1), we have

$$\langle \nabla F_{\beta_k}(x_k), \tilde{s}_k - x_k \rangle \leq \langle \nabla F_{\beta_k}(x_k), s_k - x_k \rangle + \frac{\eta_k}{2} D^2_{\mathcal{X}}(L_f + \frac{\|A\|^2}{\beta_k})$$

$$\leq \langle \nabla F_{\beta_k}(x_k), x^* - x_k \rangle + \frac{\eta_k}{2} D^2_{\mathcal{X}}(L_f + \frac{\|A\|^2}{\beta_k})$$

$$= \langle \nabla f(x_k), x^* - x_k \rangle + \langle A^T \nabla g_{\beta_k}(Ax_k), x^* - x_k \rangle + \frac{\eta_k}{2} D^2_{\mathcal{X}}(L_f + \frac{\|A\|^2}{\beta_k}),$$

where the second line follows since $s_k$ is a solution of $\min_{x \in \mathcal{X}} \langle \nabla F_{\beta_k}(x_k), x \rangle$.

Now, convexity of $f$ ensures $\langle \nabla f(x_k), x^* - x_k \rangle \leq f(x^*) - f(x_k)$. Using property (7.2), we have

$$\langle A^T \nabla g_{\beta_k}(Ax_k), x^* - x_k \rangle = \langle \nabla g_{\beta_k}(Ax_k), Ax^* - Ax_k \rangle$$

$$\leq g(Ax^*) - g_{\beta_k}(Ax_k) - \frac{\beta_k}{2} \|y_{\beta_k}(Ax_k)\|^2.$$

Putting these altogether, we get the following bound

$$F_{\beta_k}(x_{k+1}) \leq F_{\beta_k}(x_k) + \eta_k \left( f(x^*) - f(x_k) + g(Ax^*) - g_{\beta_k}(Ax_k) - \frac{\beta_k}{2} \|\nabla y_{\beta_k}(Ax_k)\|^2 \right)$$

$$+ \frac{\eta_k^2}{2} D^2_{\mathcal{X}}(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta)$$

$$= (1 - \eta_k) F_{\beta_k}(x_k) + \eta_k F(x^*) - \frac{\eta_k \beta_k}{2} \|\nabla y_{\beta_k}(Ax_k)\|^2 + \frac{\eta_k^2}{2} D^2_{\mathcal{X}}(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta). \tag{7.6}$$

Now, using (7.3), we get

$$F_{\beta_k}(x_k) = f(x_k) + g_{\beta_k}(Ax_k)$$

$$\leq f(x_k) + g_{\beta_k-1}(Ax_k) + \frac{\beta_{k-1} - \beta_k}{2} \|y_{\beta_k}(Ax_k)\|^2$$

$$= F_{\beta_k-1}(x_k) + \frac{\beta_{k-1} - \beta_k}{2} \|y_{\beta_k}(Ax_k)\|^2.$$

We combine this with (7.6) and subtract $F(x^*)$ from both sides to get

$$F_{\beta_k}(x_{k+1}) - F(x^*) \leq (1 - \eta_k)(F_{\beta_k-1}(x_k) - F(x^*)) + \frac{\eta_k^2}{2} D^2_{\mathcal{X}}(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta)$$

$$+ ((1 - \eta_k)(\beta_{k-1} - \beta_k) - \eta_k \beta_k) \frac{1}{2} \|y_{\beta_k}(Ax_k)\|^2.$$
Let us choose \( \eta_k \) and \( \beta_k \) in a way to vanish the last term. By choosing \( \eta_k = \frac{2}{k+1} \) and \( \beta_k = \frac{\beta_0}{k+1} \) for \( k \geq 1 \) with some \( \beta_0 > 0 \), we get \( (1 - \eta_k)(\beta_{k-1} - \beta_k) - \eta_k \beta_k < 0 \). Hence, we end up with

\[
F_{\beta_k}(x_{k+1}) - F(x^*) \leq (1 - \eta_k)(F_{\beta_{k-1}}(x_k) - F(x^*)) + \frac{\eta_k^2}{2} D_f^2(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta).
\]

By recursively applying this inequality, we get

\[
F_{\beta_k}(x_{k+1}) - F(x^*) \leq \prod_{j=1}^k (1 - \eta_j) \left( F_{\beta_{j-1}}(Ax_k) - F(x^*) \right) + \frac{1}{2} D_f^2(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta) \sum_{j=1}^k (1 - \eta_j)(L_f + \frac{\|A\|^2}{\beta_j}) \leq \prod_{j=1}^k (1 - \eta_j) \left( F_{\beta_{j-1}}(Ax_k) - F(x^*) \right) + \frac{1}{2} D_f^2(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta) \sum_{j=1}^k (1 - \eta_j) \leq \frac{1}{2} D_f^2(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta) \sum_{j=1}^k (1 - \eta_j),
\]

where the second line follows since \( \beta_k \leq \beta_j \) for any positive integer \( j \leq k \), and the third line since \( \eta_1 = 1 \).

Now, we use the following relation

\[
\sum_{j=1}^k \eta_j^2 \prod_{j=1}^k (1 - \eta_j) = \sum_{j=1}^k \frac{4}{(j + 1)^2} \prod_{j=1}^{k - 1} \frac{j - 1}{j + 1} = \sum_{j=1}^k \frac{4}{(j + 1)^2} \frac{(k - j)(k - j - 1)}{k(k + 1)} \leq \frac{4}{k + 1},
\]

which yields the first result of Theorem 4.1 as

\[
F_{\beta_k}(x_{k+1}) - F(x^*) \leq \frac{2}{k + 1} D_f^2(L_f + \frac{\|A\|^2}{\beta_k})(1 + \delta) = 2D_f^2(\frac{L_f}{k + 1} + \frac{\|A\|^2}{\beta_0 \sqrt{k + 1}})(1 + \delta).
\]

**Proof of Theorem 4.2**

Now, we further assume that \( g : \mathbb{R}^d \rightarrow \mathbb{R} \cup \{+\infty\} \) is \( L_g \)-Lipschitz continuous. From (7.4), we get

\[
g(Ax_{k+1}) \leq g_{\beta_k}(Ax_{k+1}) + \frac{\beta_k L_g^2}{2} = g_{\beta_k}(Ax_{k+1}) + \frac{\beta_0 L_g^2}{2\sqrt{k + 1}}.
\]

We complete the proof by adding \( f(x_{k+1}) - F(x^*) \) to both sides:

\[
F(x_{k+1}) - F(x^*) \leq F_{\beta_k}(x_{k+1}) - F(x^*) + \frac{\beta_0 L_g^2}{2\sqrt{k + 1}}.
\]

**Proof of Theorem 4.3**

From the Lagrange saddle point theory, we know that the following bound holds \( \forall x \in \mathcal{X} \) and \( \forall r \in \mathcal{K} \):

\[
f^* \leq \mathcal{L}(x, r, y^*) = f(x) + \langle y^*, Ax - r \rangle \leq f(x) + \|y^*\|\|Ax - r\|,
\]

Since \( x_{k+1} \in \mathcal{X} \), we get

\[
f(x_{k+1}) - f^* \geq - \min_{r \in \mathcal{K}} \|y^*\|\|Ax_{k+1} - r\| = -\|y^*\|\text{dist}(Ax_{k+1}, \mathcal{K}). \tag{7.7}
\]
This proves the first bound in Theorem 4.3.

The second bound directly follows by Theorem 4.1 as

\[
\frac{f(x_{k+1}) - f^*}{\frac{1}{2\beta_k} \text{dist}^2(Ax_{k+1}, \mathcal{K})} \leq \frac{2D_k^2}{\beta_0 \sqrt{k+1}} \left( L_f + \frac{2\beta_k}{\beta_0} \right) (1 + \delta).
\]

Now, we combine this with (7.7), and we get

\[
-\|y^*\| \text{dist}(Ax_{k+1}, \mathcal{K}) + \frac{1}{2\beta_k} \text{dist}^2(Ax_{k+1}, \mathcal{K}) \leq 2D_k^2 \left( L_f + \frac{2\beta_k}{\beta_0} \right) (1 + \delta)
\leq 2D_k^2 \frac{\beta_k}{\beta_0} (L_f + \frac{2\beta_k}{\beta_0}) (1 + \delta).
\]

This is a second order inequality in terms of dist(Ax_k, \mathcal{K}). Solving this inequality, we get

\[
\text{dist}(Ax_{k+1}, \mathcal{K}) \leq \beta_k \left( \|y^*\| + \sqrt{\|y^*\|^2 + 4D_k^2 \frac{1}{\beta_0} \left( L_f + \frac{2\beta_k}{\beta_0} \right) \right) (1 + \delta)
\]

\[
\leq \frac{2\beta_0}{\sqrt{k+1}} \left( \|y^*\| + D_k \sqrt{\frac{1}{\beta_0} \left( L_f + \frac{2\beta_k}{\beta_0} \right) \right) (1 + \delta).
\]

**Proof of Theorem 4.4**

Let us define the multiplicative error \( \delta \) of the LMO:

\[
\langle v_k, \tilde{s}_k - x_k \rangle \leq \delta \langle v_k, s_k - x_k \rangle
\]

(7.8)

For the proof we assume that \( x_1 \) is feasible. First, we use the smoothness of \( F_{\beta_k} \) to upper bound the progress. Note that \( F_{\beta_k} \) is \( (L_f + A^2/\beta_k) \)-smooth.

\[
F_{\beta_k}(x_{k+1}) \leq F_{\beta_k}(x_k) + \eta_k \langle \nabla F_{\beta_k}(x_k), \tilde{s}_k - x_k \rangle + \frac{\eta_k^2}{2} \|\tilde{s}_k - x_k\|^2 (L_f + A^2/\beta_k)
\]

\[
\leq F_{\beta_k}(x_k) + \eta_k \langle \nabla F_{\beta_k}(x_k), \tilde{s}_k - x_k \rangle + \frac{\eta_k^2}{2} D_k^2 (L_f + \frac{A^2}{\beta_k}),
\]

(7.9)

where \( \tilde{s}_k \) denotes the atom selected by the inexact linear minimization oracle, and the second inequality follows since \( \tilde{s}_k \in \mathcal{X} \).

By definition of inexact oracle (7.8), we have

\[
\langle \nabla F_{\beta_k}(x_k), \tilde{s}_k - x_k \rangle \leq \delta \langle \nabla F_{\beta_k}(x_k), s_k - x_k \rangle
\leq \delta \langle \nabla F_{\beta_k}(x_k), x^* - x_k \rangle
= \delta \langle \nabla f(x_k), x^* - x_k \rangle + \delta \langle A^\top \nabla g_{\beta_k}(Ax_k), x^* - x_k \rangle,
\]

where the second line follows since \( s_k \) is a solution of \( \min_{x \in \mathcal{X}} \langle \nabla F_{\beta_k}(x_k), x \rangle \).

Now, convexity of \( f \) ensures \( \langle \nabla f(x_k), x^* - x_k \rangle \leq f(x^*) - f(x_k) \). Using property (7.2), we have

\[
\langle A^\top \nabla g_{\beta_k}(Ax_k), x^* - x_k \rangle = \langle \nabla g_{\beta_k}(Ax_k), Ax^* - Ax_k \rangle
\]

\[
\leq g(Ax^*) - g_{\beta_k}(Ax_k) - \frac{\beta_k}{2} \|y^*_{\beta_k}(Ax_k)\|^2.
\]

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Putting these altogether, we get the following bound
\[
F_{\beta_k}(x_{k+1}) \leq F_{\beta_k}(x_k) + \eta_k \delta \left( f(x^*) - f(x_k) + g(Ax^*) - g_{\beta_k}(Ax_k) - \frac{\beta_k}{2} \|\nabla y_{\beta_k}(Ax_k)\|^2 \right) \\
+ \frac{\eta_k^2}{2} D_X^2(L_f + \frac{\|A\|^2}{\beta_k}) \\
= (1 - \delta \eta_k) F_{\beta_k}(x_k) + \delta \eta_k \frac{\beta_k - \beta_k - \beta_k}{2} \|y_{\beta_k}(Ax_k)\|^2 \\
+ \frac{\eta_k^2}{2} D_X^2(L_f + \frac{\|A\|^2}{\beta_k}).
\]
(7.10)

Now, using (7.3), we get
\[
F_{\beta_k}(x_k) = f(x_k) + g_{\beta_k}(Ax_k) \\
\leq f(x_k) + g_{\beta_{k-1}}(Ax_k) + \frac{\beta_{k-1} - \beta_k}{2} \|y_{\beta_k}(Ax_k)\|^2 \\
= F_{\beta_{k-1}}(x_k) + \frac{\beta_{k-1} - \beta_k}{2} \|y_{\beta_k}(Ax_k)\|^2.
\]

We combine this with (7.10) and subtract \(F(x^*)\) from both sides to get
\[
F_{\beta_k}(x_{k+1}) - F(x^*) \leq (1 - \delta \eta_k)(F_{\beta_{k-1}}(x_k) - F(x^*)) + \frac{\eta_k^2}{2} D_X^2(L_f + \frac{\|A\|^2}{\beta_k}) \\
+ ((1 - \delta \eta_k)(\beta_{k-1} - \beta_k) - \delta \eta_k \beta_k) \frac{1}{2} \|\nabla y_{\beta_k}(Ax_k)\|^2.
\]

By choosing \(\eta_k = \frac{2}{\delta(k-1)+2}\) and \(\beta_k = \frac{\beta_0}{\sqrt{\delta k+1}}\) for some \(\beta_0 > 0\), we get \((1 - \delta \eta_k)(\beta_{k-1} - \beta_k) - \delta \eta_k \beta_k < 0\) for any \(k \geq 1\), hence we end up with
\[
F_{\beta_k}(x_{k+1}) - F(x^*) \leq (1 - \delta \eta_k)(F_{\beta_{k-1}}(x_k) - F(x^*)) + \frac{\eta_k^2}{2} D_X^2(L_f + \frac{\|A\|^2}{\beta_k}).
\]

Let us call for simplicity \(C := D_X^2(L_f + \frac{\|A\|^2}{\beta_k})\), \(E_{k+1} := F_{\beta_k}(x_{k+1}) - F(x^*)\) Therefore, we have
\[
E_{k+1} \leq (1 - \delta \eta_k)E_k + \frac{\eta_k^2}{2} C
\]
(7.11)

We now show by induction that:
\[
E_k \leq 2 \frac{1}{2} + \frac{1}{\delta(k-1) + 2}
\]

The base case \(k = 1\) is trivial as \(C > 0\). Call for simplicity \(K := \delta(k-1) + 2\). Note that \(K \geq 2\).

Under this notation we can write \(\eta_k = \frac{2}{\delta(k-1)+2} = \frac{2}{K}\) For the induction step, we add a positive term
\( (E_1 \text{ is positive as } x_1 \text{ is assumed feasible}) \) to (7.11) and use the induction hypothesis:

\[
E_{k+1} \leq (1 - \delta \eta_k)E_k + \frac{\eta_k^2}{2} C + 2\delta \frac{E_1}{K^2}
\]

\[
\leq (1 - \delta \frac{2}{K})E_k + \frac{2}{K^2} C + 2\delta \frac{E_1}{K^2}
\]

\[
\leq (1 - \delta \frac{2}{K})^2 \frac{\frac{1}{2}C + E_1}{K} + \frac{2}{K^2} C + 2\delta \frac{E_1}{K^2}
\]

\[
= (1 - \delta \frac{2}{K})^2 \frac{\frac{1}{2}C + E_1}{K} + 2\delta \left( \frac{\frac{1}{2}C}{K^2} + \frac{E_1}{K^2} \right)
\]

\[
= \frac{\frac{1}{2}C + E_1}{K} \left( 1 - \delta \frac{2}{K} + \delta \right)
\]

\[
= \frac{\frac{1}{2}C + E_1}{K} \left( 1 - \delta \frac{\delta}{K} \right)
\]

\[
\leq \frac{\frac{1}{2}C + E_1}{K + \delta}
\]

noting that \( K + \delta = \delta k + 2 \) concludes the proof.

Proof of Theorems 4.5 and 4.6 follows similarly to the proofs of Theorems 4.2 and 4.3.